

Homework problems:

1. Consider the following context-free grammars:

$$\begin{array}{ll} \text{(a)} & A \rightarrow aAcc \mid B \\ & B \rightarrow bBc \mid \varepsilon \end{array} \qquad \text{(b)} \quad S \rightarrow +S- \mid SS \mid \varepsilon$$

Give a derivation for the sentence $abccc$ according to grammar (a), and a derivation for the sentence $++-+- -+-$ according to grammar (b). Describe the language generated by each grammar verbally as simply as you can.

2. A *palindrome* is a string w such that $w = w^R$. (E.g. “MADAMIMADAM”, “ABLEWASIEREISAWELBA,” cf. <http://www.palindromes.org/>.) Consider the set of palindromes over the alphabet $\{a, b\}$:

$$\text{PAL} = \{w \in \{a, b\}^* \mid w = w^R\}.$$

Design a context-free grammar generating the language. (*Hint*: Note that a string $w \in \text{PAL}$, if and only if it is of the form $w = uXu^R$, where $X = a, b$ or ε .)

3. Consider the following grammar generating a certain type of list structures:

$$S \rightarrow (S) \mid S, S \mid a.$$

- (a) Based on the above grammar, give a leftmost and rightmost derivation and a parse tree for the sentence “ $(a, (a))$ ”.
- (b) Prove that the grammar is ambiguous.
- (c) Design an unambiguous grammar generating the same language.

Demonstration problems:

4. Prove that the class of context-free languages is closed under unions, concatenations, and the Kleene star operation, i.e. if the languages $L_1, L_2 \subseteq \Sigma^*$ are context-free, then so are the languages $L_1 \cup L_2$, L_1L_2 and L_1^* .
5. Design a context-free grammar describing the syntax of simple “programs” of the following form: a program consists of nested **for** loops, compound statements enclosed by **begin-end** pairs and elementary operations **a**. Thus, a “program” in this language looks something like this:

```
a;  
for 3 times do  
begin  
  for 5 times do a;  
  a; a  
end.
```

For simplicity, you may assume that the loop counters are always integer constants in the range $0, \dots, 9$.

6. (a) Prove that the following context-free grammar is ambiguous:

$$\begin{array}{l} S \rightarrow \mathbf{if\ } b \mathbf{\ then\ } S \\ S \rightarrow \mathbf{if\ } b \mathbf{\ then\ } S \mathbf{\ else\ } S \\ S \rightarrow s. \end{array}$$

- (b) Design an unambiguous grammar that is equivalent to the grammar in item (a), i.e. that generates the same language. (*Hint*: Introduce new nonterminals B and U that generate, respectively, only “balanced” and “unbalanced” **if-then-else**-sequences.)