

**Ordinary exercises:**

1. Construct a two-tape Turing machine that decides the language:

$$\{a^n b^n c^n \mid n \geq 0\}.$$

Use both tapes during the computation.

Convention: the input is read and the answer is written on the first tape as in the case of a single-tape machine. The second tape is initially empty and its read/write head is positioned in the first tape position. It doesn't matter what the second tape contains in the end of the computation.

2. Form an unrestricted grammar (=type 0 grammar)  $G$  such that:

$$L(G) = \{ww \mid w \in \{a, b\}^*\}.$$

3. It can be proved that every unrestricted grammar  $G$  can be converted into an equivalent grammar  $G'$  such that all rules of  $G'$  are of the form:  $uAv \rightarrow uvw$  where  $A \in V - \Sigma$  and  $u, v, w \in V^*$ .

Do the conversion to the grammar  $G = (V, \Sigma, R, S)$ , where

$$\begin{aligned}\Sigma &= \{a\}, \\ V &= \Sigma \cup \{S, [, ], A, N\}, \text{ and} \\ R &= \{S \rightarrow [NA], S \rightarrow a, [N \rightarrow [NN, NA \rightarrow AAN, N] \rightarrow], [A \rightarrow a[, [] \rightarrow e\}.\end{aligned}$$

Hint: Some of the rules are already in the desired form. You can convert the other rules by adding new non-terminals and splitting a rule into two or more new rules.

**Demonstration exercises:**

4. Construct an unrestricted grammar that generates the language:

$$L = \{a^{n^2} \mid n \geq 0\}$$

5. Prove that the following problem is undecidable:

Let  $M$  be a Turing machine. Does  $M$  stop when it is given the empty string  $e$  as input.

6. Show that the class of Turing acceptable languages is closed under union and intersection.