

This is the last set of tutorial problems. The course exam is on Wednesday 8 May, 9–12 a.m. Remember to register for the exam via TOPI, and fill in and return the course feedback form. (Instructions for the latter will be available on the course's WWW page on Tuesday 23 April.)

Homework problems:

1. You are being offered the following programming assignment:

Intel Septium code optimisation

A large company producing embedded systems software would like to have a code optimiser that, given as input any machine language program for the new Intel Septium processor, will output the smallest machine language program that is functionally equivalent to the given one (i.e. has the same input-output behaviour).

Your comments? Under what conditions would you accept the assignment? Justify your answer.

2. Formulate Rice's general undecidability theorem and the related notions ("decidable properties of Turing machines" etc.) in terms of C programs rather than Turing machines. [If your textbook doesn't discuss this fundamental result, look it up in some reference work on computability theory or on the Internet.]
3. Consider application programs running under some given operating system. Let us say that a program P is *dangerous*, if it on some input modifies the operating system's system files. A *general purpose virus tester* is a program that receives as input an arbitrary application program text P , analyses it and returns output "DANGER", if the program is dangerous, and "OK" otherwise. Show that if any dangerous programs exist at all, then general-purpose virus testing is impossible.

Demonstration problems:

4. Prove, without appealing to Rice's theorem, that the following problem is undecidable:

Given a Turing machine M ; does M accept the empty string?

5. Prove the following connections between recursive functions and languages:

- (i) A language $A \subseteq \Sigma^*$ is recursive ("Turing-decidable"), if and only its characteristic function

$$\chi_A : \Sigma^* \rightarrow \{0, 1\}, \quad \chi_A(x) = \begin{cases} 1, & \text{if } x \in A; \\ 0, & \text{if } x \notin A \end{cases}$$

is a recursive ("Turing-computable") function.

- (ii) A language $A \subseteq \Sigma^*$ is recursively enumerable ("semidecidable", "Turing-recognisable"), if and only if either $A = \emptyset$ or there exists a recursive function $g : \{0, 1\}^* \rightarrow \Sigma^*$ such that

$$A = \{g(x) \mid x \in \{0, 1\}^*\}.$$